

Economics 201B–Second Half

Lecture 6

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The Second Welfare Theorem in the Arrow Debreu Economy

Theorem 1 (Second Welfare Theorem) (Pure Exchange Case) *If x^* is Pareto Optimal in a pure exchange economy, with strongly monotone, continuous, convex preferences, there exists a price vector p^* and an income transfer T such that (p^*, x^*, T) is a Walrasian Equilibrium with Transfers.*

Outline of Proof:

- Let

$$B_i = \{x'_i - x_i^* : x'_i \succ_i x_i^*\}$$

$$B = \sum_{i=1}^I B_i = \{b_1 + \cdots + b_I : b_i \in B_i\}$$

Then $0 \notin B$ (if it were, we'd have a Pareto improvement).

- By Minkowski's Theorem, find $p^* \neq 0$ such that

$$\inf p^* \cdot B \geq 0$$

- Show $(\mathbf{R}_+^L \setminus \{0\}) \subset B_i$ and hence $p^* \geq 0$.

- Show $\inf p^* \cdot B_i = 0$ for each i .
- Define T to make x_i^* affordable at p^* :

$$T_i = p^* \cdot x_i^* - p^* \cdot \omega_i$$

Show $\sum_{i=1}^I T_i = 0$ and

$$x_i^* \in Q_i(p^*, T)$$

- Use strong monotonicity to show that $p^* \gg 0$.
- Show

$$p^* \gg 0 \Rightarrow Q_i(p^*, T) = D_i(p^*, T)$$

Now, for the details:

- Let

$$B_i = \{x'_i - x_i^* : x'_i \succ_i x_i^*\}$$

$$B = \sum_{i=1}^I B_i = \{b_1 + \cdots + b_I : b_i \in B_i\}$$

Claim:

$$0 \notin B$$

If $0 \in B$, there exists $b_i \in B_i$ such that

$$\sum_{i=1}^I b_i = 0$$

Let

$$x'_i = x_i^* + b_i$$

Since $x'_i - x_i^* = b_i \in B_i$, we have

$$x'_i \succ_i x_i^*$$

$$\begin{aligned} \sum_{i=1}^I x'_i &= \sum_{i=1}^I (x_i^* + b_i) \\ &= \sum_{i=1}^I x_i^* + \sum_{i=1}^I b_i \\ &= \sum_{i=1}^I x_i^* \\ &= \bar{\omega} \end{aligned}$$

Therefore, x' is a feasible allocation, x' Pareto improves x^* , so x^* is not Pareto Optimal, contradiction. Therefore, $0 \notin B$.

•

$$\exists_{p^* \neq 0} \inf p^* \cdot B \geq 0$$

B_i is convex, so B is convex (easy exercise). By Minkowski's Theorem, there exists $p^* \neq 0$ such that

$$0 = p^* \cdot 0 \leq \inf p^* \cdot B = \sum_{i=1}^I \inf p^* \cdot B_i$$

The fact that $\inf p^* \cdot B = \sum_{i=1}^I \inf p^* \cdot B_i$ is an exercise; once you figure out what you have to prove, it is obvious.

• We claim that $p^* \geq 0$.

Suppose not, so $p_\ell^* < 0$ for some ℓ , WLOG $p_1^* < 0$.

Let

$$x'_i = x_i^* + \left(-\frac{1}{p_1^*}, 0, \dots, 0 \right)$$

By strong monotonicity, $x'_i \succ_i x_i^*$, so

$$\left(-\frac{1}{p_1^*}, 0, \dots, 0 \right) \in B_i$$

So

$$\begin{aligned} \inf p^* \cdot B_i &\leq p^* \cdot \left(-\frac{1}{p_1^*}, 0, \dots, 0 \right) \\ &= -1 < 0 \\ \inf p^* \cdot B &= \sum_{i=1}^I \inf p^* \cdot B_i \\ &\leq -I \\ &< 0 \end{aligned}$$

a contradiction that shows $p^* \geq 0$.

- We claim that $\inf p^* \cdot B_i = 0$ for each i :
Suppose $\varepsilon > 0$. By strong monotonicity,

$$x_i^* + (\varepsilon, \dots, \varepsilon) \succ_i x_i^*$$

so

$$(\varepsilon, \dots, \varepsilon) \in B_i$$

so

$$\inf p^* \cdot B_i \leq p^* \cdot (\varepsilon, \dots, \varepsilon)$$

Since ε is an arbitrary positive number, $\inf p^* \cdot B_i$ is less than every positive number, so

$$\inf p^* \cdot B_i \leq 0$$

Since $\sum_{i=1}^I \inf p^* \cdot B_i \geq 0$,

$$\inf p^* \cdot B_i = 0 \quad (i = 1, \dots, I)$$

- Define T to make x_i^* affordable at p^* . We claim that T is an income transfer and

$$x_i^* \in Q_i(p^*, T)$$

Let

$$T_i = p^* \cdot x_i^* - p^* \cdot \omega_i$$

$$\begin{aligned} \sum_{i=1}^I T_i &= \sum_{i=1}^I (p^* \cdot x_i^* - p^* \cdot \omega_i) \\ &= p^* \cdot \left(\sum_{i=1}^I x_i^* - \sum_{i=1}^I \omega_i \right) \\ &= p^* \cdot (\bar{\omega} - \bar{\omega}) \\ &= 0 \end{aligned}$$

so T is an income transfer.

$$\begin{aligned} p^* \cdot x_i^* &= p^* \cdot (\omega_i + (x_i^* - \omega_i)) \\ &= p^* \cdot \omega_i + p^* \cdot (x_i^* - \omega_i) \\ &= p^* \cdot \omega_i + T_i \end{aligned}$$

so

$$x_i^* \in B_i(p^*, T)$$

If $x'_i \succ_i x_i^*$, then $x'_i - x_i^* \in B_i$, so

$$\begin{aligned} p^* \cdot x'_i &= p^* \cdot (x_i^* + (x'_i - x_i^*)) \\ &= p^* \cdot x_i^* + p^* \cdot (x'_i - x_i^*) \\ &\geq p^* \cdot x_i^* + \inf p^* \cdot B_i \\ &= p^* \cdot x_i^* \\ &= p^* \cdot \omega_i + T_i \end{aligned}$$

so

$$x_i^* \in Q(p^*, T)$$

- Use strong monotonicity to show that $p^* \gg 0$.

Lemma 2 *If \succeq_i is continuous and complete, and $x \succ_i y$, then there exists $\varepsilon > 0$ such that*

$$(B(x, \varepsilon) \cap X_i) \succ_i y$$

Proof: If not, we can find $x_n \rightarrow x$, $x_n \in X_i$, $x_n \not\succeq_i y$; by completeness, we have $y \succeq_i x_n$ for each n . Since \succeq_i is continuous, $y \succeq_i x$, so $x \not\succeq_i y$, a contradiction which proves the lemma. ■

Since $p^* \geq 0$ and $p^* \neq 0$, $p^* > 0$; since in addition $\bar{\omega} \gg 0$, $p^* \cdot \bar{\omega} > 0$, so

$$p^* \cdot \omega_i + T_i > 0 \text{ for some } i$$

If $p_\ell^* = 0$ for some ℓ (WLOG $\ell = 1$), let

$$x'_i = x_i^* + (1, 0, \dots, 0)$$

By strong monotonicity, $x'_i \succ_i x_i^*$.

$$p^* \cdot x'_i = p^* \cdot x_i^* = p^* \cdot \omega_i + T_i > 0$$

Find ℓ (WLOG $\ell = 2$) such that

$$p_\ell^* > 0, \quad x'_{2i} > 0$$

Since $x'_i \succ_i x_i^*$, let $\varepsilon > 0$ be chosen to satisfy the conclusion of the Lemma. If necessary, we may make ε smaller to ensure that $\varepsilon \leq 2x'_{2i}$. Let

$$x''_i = x'_i - (0, \varepsilon/2, 0, \dots, 0)$$

Since $X_i = \mathbf{R}_+^L$, $x''_i \in X_i$, so by the Lemma, $x''_i \succ_i x_i^*$. But $p^* \cdot x''_i < p \cdot x'_i = p^* \cdot \omega_i + T_i$, which shows that $x_i^* \notin Q_i(p^*, T)$, a contradiction which proves that $p^* \gg 0$.

• Show

$$p^* \gg 0 \Rightarrow Q_i(p^*, T) = D_i(p^*, T)$$

– Case 1: $p^* \cdot \omega_i + T_i = 0$. Since $p^* \gg 0$, $B_i(p^*, T) = \{0\}$, so

$$Q_i(p^*, T) = D_i(p^*, T) = \{0\}$$

– Case 2: $p^* \cdot \omega_i + T_i > 0$

☞ Suppose $x \in Q_i(p^*, T)$ but $x \notin D_i(p^*, T)$. Then there exists $z \succ_i x$ such that $z \in B_i(p^*, T)$, hence $p^* \cdot z \leq p^* \cdot \omega_i + T_i$. Since $x \in Q_i(p^*, T)$, $p^* \cdot z \geq p^* \cdot \omega_i + T_i$, so

$$p^* \cdot z = p^* \cdot \omega_i + T_i > 0$$

By Lemma 2, there exists $\varepsilon > 0$ such that

$$|z' - z| < \varepsilon, z' \in \mathbf{R}_+^L \Rightarrow z' \succ x$$

Let

$$z' = z \left(1 - \frac{\varepsilon}{2|z|} \right)$$

Since $z \in \mathbf{R}_+^L$, $z' \in \mathbf{R}_+^L$.

$$|z' - z| = \left| \frac{\varepsilon z}{2|z|} \right| = \frac{\varepsilon}{2} < \varepsilon$$

so $z' \succ x$.

$$\begin{aligned} p^* \cdot z' &= p^* \cdot z \left(1 - \frac{\varepsilon}{2|z|} \right) \\ &= (p^* \cdot \omega_i + T_i) \left(1 - \frac{\varepsilon}{2|z|} \right) \\ &< p^* \cdot \omega_i + T_i \end{aligned}$$

which contradicts the assumption that $x \in Q_i(p^*, T)$.

This shows $Q_i(p^*, T) \subset D_i(p^*, T)$; since clearly $D_i(p^*, T) \subset Q_i(p^*, T)$, $Q_i(p^*, T) = D_i(p^*, T)$.

What if preferences are not convex?

- Second Welfare Theorem may fail if preferences are nonconvex.
- Diagram gives an economy with two goods and two agents, and a Pareto optimum x^* so that so that the utility levels of x^* cannot be approximated by a Walrasian Equilibrium with Transfers.
- If p^* is the price which locally supports x^* , and T is the income transfer which makes x affordable with respect to the prices p^* , there is a unique Walrasian equilibrium with transfers (z^*, q^*, T) ; z^* is much more favorable to agent I and much less favorable to agent II than x^* is.
- This is the worst that can happen under standard assumptions on preferences. Given a Pareto optimum x^* , there is a Walrasian quasiequilibrium with transfers (z^*, p^*, T) such that all but L people are indifferent between x^* and z^* . Those L people are treated quite harshly (they get zero consumption). One could be less harsh and give these L people carefully chosen consumption bundles in the convex hull of their quasidemand sets, *but one would then have to forbid them from trading*, a prohibition

that would in practice be difficult to enforce.