

# **Microeconomics III**

**Nash equilibrium II**  
**(Apr 15, 2012)**

**School of Economics**  
**The Interdisciplinary Center (IDC), Herzliya**

## Randomization

Recall that a strategic game is a triple  $\langle N, (A_i), (\succsim_i) \rangle$  where

- $N$  is a finite set of players, and for each player  $i \in N$
- a non-empty set  $A_i$  of actions
- a preference relation  $\succsim_i$  on the set  $A = \times_{j \in N} A_j$  of possible outcomes.

or a triple  $\langle N, (A_i), (u_i) \rangle$  when  $\succsim_i$  can be represented by a utility function  $u_i : A \rightarrow \mathbb{R}$ .

Suppose that,

- each player  $i$  can randomize among all her strategies so choices are not deterministic, and
- player  $i$ 's preferences over lotteries on  $A$  can be represented by  $vNM$  expected utility function.

Then, we need to add these specifications to the primitives of the model of strategic game  $\langle N, (A_i), (\succsim_i) \rangle$ .

A mixed strategy of player  $i$  is  $\alpha_i \in \Delta(A_i)$  where  $\Delta(A_i)$  is the set of all probability distributions over  $A_i$ .

- A profile  $(\alpha_i)_{i \in N}$  of mixed strategies induces a probability distribution over the set  $A$ .
- Assuming independence, the probability of an action profile (outcome)  $a$  is then

$$\prod_{i \in N} \alpha_i(a_i).$$

A *vNM* utility function

$$U_i : \times_{j \in N} \Delta(A_j) \rightarrow \mathbb{R}$$

represents player *i*'s preferences over the set of lotteries over *A*.

The mixed extension of a the strategic game  $\langle N, (A_i), (u_i) \rangle$  is the strategic game

$$\langle N, (\Delta(A_i)), (U_i) \rangle .$$

## Preferences toward risk

The standard model of decisions under risk (known probabilities) is based on von Neumann and Morgenstern Expected Utility Theory.

Consider a set of *lotteries*, or gambles, (outcomes and probabilities). A fundamental axiom about preferences toward risk is *independence*:

For any lotteries  $x, y, z$  and  $0 < \alpha < 1$

$$x \succ y \text{ implies } \alpha x + (1 - \alpha)z \succ \alpha y + (1 - \alpha)z.$$

Expected Utility Theory has some very convenient properties for analyzing choice under uncertainty.

To clarify, we will consider the *utility* that a consumer gets from her or his income.

More precisely, from the consumption bundle that the consumer's income can buy.

## Behavioral economics

### Allais (1953) I

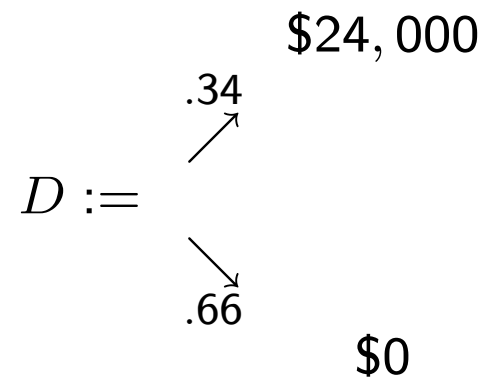
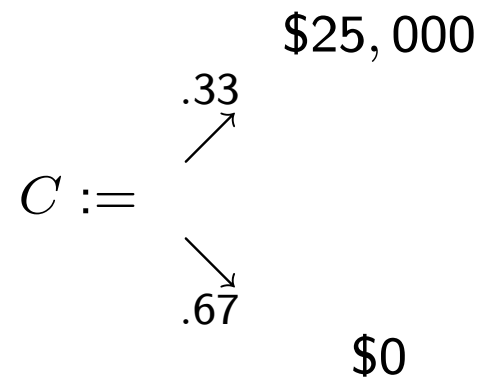
- Choose between the two gambles:

$$A := \begin{array}{l} \begin{array}{c} \nearrow .33 \\ \longrightarrow .66 \\ \searrow .01 \end{array} \begin{array}{l} \$25,000 \\ \$24,000 \\ \$0 \end{array} \end{array} \quad B := \xrightarrow{1} \$24,000$$



## Allais (1953) II

- Choose between the two gambles:



## Two results on mixed strategy Nash equilibrium

Let  $G = \langle N, (A_i), (u_i) \rangle$  be a strategic game and  $G' = \langle N, (\Delta(A_i)), (U_i) \rangle$  be its mixed extension.

[1] If  $a \in NE(G)$  then  $a \in NE(G')$ .

[2]  $\alpha \in NE(G')$  if and only if

$$U_i(\alpha_{-i}, a_i) \geq U_i(\alpha_{-i}, a'_i)$$

for all  $a'_i$  and all  $\alpha_i(a_i) > 0$ .

[1] Proof: If  $a \in NE(G)$  then

$$u_i(a_{-i}, a_i) \geq u_i(a_{-i}, a'_i) \quad \forall i \in N \text{ and } \forall a'_i \in A_i.$$

Then, by the linearity of  $U_i$  in  $\alpha_i$

$$U_i(a_{-i}, a_i) \geq U_i(a_{-i}, \alpha_i) \quad \forall i \in N \text{ and } \forall \alpha_i \in \Delta(A_i)$$

and thus  $a \in NE(G')$ .

[2] Proof: Let  $\alpha \in NE(G')$

Suppose that  $\exists a_i \in A_i$  such that  $\alpha_i(a_i) > 0$  and

$$U_i(\alpha_{-i}, a'_i) \geq U_i(\alpha_{-i}, a_i) \text{ for some } a'_i \neq a_i.$$

Then, player  $i$  can increase her payoff by transferring probability from  $a_i$  to  $a'_i$  so  $\alpha$  is not a  $NE$ .

This implies that  $U_i(\alpha_{-i}, a_i) = U_i(\alpha_{-i}, a'_i)$  for all  $a_i, a'_i$  in the support of  $\alpha$ .

## Evolutionary stability

A single population of players. Players interact with each other pair-wise and randomly matched.

Players are assigned modes of behavior (mutation). Utility measures each player's ability to survive.

$\varepsilon$  of players consists of mutants taking action  $a$  while others take action  $a^*$ .

## Evolutionary stable strategy (*ESS*)

Consider a payoff symmetric game  $G = \langle \{1, 2\}, (A, A), (u_i) \rangle$  where  $u_1(a) = u_2(a')$  when  $a'$  is obtained from  $a$  by exchanging  $a_1$  and  $a_2$ .

$a^* \in A$  is *ESS* iff for any  $a \in A$ ,  $a \neq a^*$  and  $\varepsilon > 0$  sufficiently small

$$(1 - \varepsilon)u(a^*, a^*) + \varepsilon u(a^*, a) > (1 - \varepsilon)u(a, a^*) + \varepsilon u(a, a)$$

which is satisfied iff for any  $a \neq a^*$  either

$$u(a^*, a^*) > u(a, a^*)$$

or

$$u(a^*, a^*) = u(a, a^*) \text{ and } u(a^*, a) > u(a, a)$$

## Three results on *ESS*

[1] If  $a^*$  is an *ESS* then  $(a^*, a^*)$  is a *NE*.

Suppose not. Then, there exists a strategy  $a \in A$  such that

$$u(a, a^*) > u(a^*, a^*).$$

But, for  $\varepsilon$  small enough

$$(1 - \varepsilon)u(a^*, a^*) + \varepsilon u(a^*, a) < (1 - \varepsilon)u(a, a^*) + \varepsilon u(a, a)$$

and thus  $a^*$  is not an *ESS*.

[2] If  $(a^*, a^*)$  is a strict  $NE$  ( $u(a^*, a^*) > u(a, a^*)$  for all  $a \in A$ ) then  $a^*$  is an  $ESS$ .

Suppose  $a^*$  is not an  $ESS$ . Then either

$$u(a^*, a^*) \leq u(a, a^*)$$

or

$$u(a^*, a^*) = u(a, a^*) \text{ and } u(a^*, a) \leq u(a, a).$$

so  $(a^*, a^*)$  can be a  $NE$  but not a strict  $NE$ .



[3] A  $2 \times 2$  game  $G = \langle \{1, 2\}, (A, A), (u_i) \rangle$  where  $u_i(a) \neq u_i(a')$  for any  $a, a'$  has a mixed strategy which is *ESS*

	$a$	$a'$
$a$	$w, w$	$x, y$
$a'$	$y, x$	$z, z$

If  $w > y$  or  $z > x$  then  $(a, a)$  or  $(a', a')$  are strict *NE*, and thus  $a$  or  $a'$  are *ESS*.

If  $w < y$  and  $z < x$  then there is a unique symmetric mixed strategy *NE*  $(\alpha^*, \alpha^*)$  where

$$\alpha^*(a) = (z - x) / (w - y + z - x)$$

and  $u(\alpha^*, \alpha) > u(\alpha, \alpha)$  for any  $\alpha \neq \alpha^*$ .

## Strictly competitive games

A strategic game  $\langle \{1, 2\}, (A_i), (\succsim_i) \rangle$  is strictly competitive if for any  $a \in A$  and  $b \in A$  we have  $a \succsim_1 b$  if and only if  $b \succsim_2 a$ .

	<i>L</i>	<i>R</i>
<i>T</i>	$A, -A$	$B, -B$
<i>B</i>	$C, -C$	$D, -D$

## Maxminimization

A max min mixed strategy of player  $i$  is a mixed strategy that solves the problem

$$\max_{\alpha_i \in \Delta A_i} \min_{\alpha_{-i} \in \Delta A_{-i}} U_i(\alpha_i, \alpha_{-i})$$

A player's payoff in  $\alpha^* \in NE(G)$  is at least her max min payoff:

$$U_i(\alpha^*) \geq U_i(\alpha_i, \alpha_{-i}^*) \geq \min_{\alpha_{-i} \in \Delta A_{-i}} U_i(\alpha_i, \alpha_{-i})$$

and thus

$$U_i(\alpha^*) \geq \max_{\alpha_i \in \Delta A_i} \min_{\alpha_{-i} \in \Delta A_{-i}} U_i(\alpha_i, \alpha_{-i})$$

since the above holds for all  $\alpha_i \in \Delta(A_i)$ .

## Two min-max results

$$[1] \quad \max_{\alpha_i \in \Delta A_i} \min_{\alpha_{-i} \in \Delta A_{-i}} U_i(\alpha_i, \alpha_{-i}) \leq \min_{\alpha_{-i} \in \Delta A_{-i}} \max_{\alpha_i \in \Delta A_i} U_i(\alpha_i, \alpha_{-i})$$

For every  $\alpha'_i$

$$\min_{\alpha_{-i}} U_i(\alpha'_i, \alpha_{-i}) \leq U_i(\alpha'_i, \alpha'_{-i})$$

and thus

$$\min_{\alpha_{-i}} U_i(\alpha'_i, \alpha_{-i}) \leq \max_{\alpha_i} U_i(\alpha_i, \alpha'_{-i})$$

However, since the above holds for every  $\alpha'_i$  and  $\alpha'_{-i}$  it must hold for the “best” and “worst” such choices

$$\max_{\alpha_i} \min_{\alpha_{-i}} U_i(\alpha_i, \alpha_{-i}) \leq \min_{\alpha_{-i}} \max_{\alpha_i} U_i(\alpha_i, \alpha_{-i}).$$

[2] In a zero-sum game

$$\max_{\alpha_1 \in \Delta A_1} \min_{\alpha_2 \in \Delta A_2} U_1(\alpha_1, \alpha_2) = \min_{\alpha_2 \in \Delta A_2} \max_{\alpha_1 \in \Delta A_1} U_1(\alpha_1, \alpha_2) = U_1(\alpha^*)$$

$\Leftarrow$  Since  $\alpha^* \in NE(G)$

$$U_1(\alpha^*) = \max_{\alpha_1 \in \Delta A_1} U_1(\alpha_1, \alpha_2^*) \geq \min_{\alpha_2 \in \Delta A_2} \max_{\alpha_1 \in \Delta A_1} U_1(\alpha_1, \alpha_2)$$

and since  $U_1 = -U_2$  at the same time

$$U_1(\alpha^*) = \min_{\alpha_2 \in \Delta A_2} U_1(\alpha_1^*, \alpha_2) \leq \max_{\alpha_1 \in \Delta A_1} \min_{\alpha_2 \in \Delta A_2} U_1(\alpha_1, \alpha_2)$$

Hence,

$$\max_{\alpha_1 \in \Delta A_1} \min_{\alpha_2 \in \Delta A_2} U_1(\alpha_1, \alpha_2) \geq \min_{\alpha_2 \in \Delta A_2} \max_{\alpha_1 \in \Delta A_1} U_1(\alpha_1, \alpha_2)$$

which together with [1] gives the desired conclusion.

$\Rightarrow$  Let  $\alpha_1^{\max}$  be player 1's max min strategy and  $\alpha_2^{\min}$  be player 2's min max strategy. Then,

$$\begin{aligned} \max_{\alpha_1 \in \Delta A_1} \min_{\alpha_2 \in \Delta A_2} U_1(\alpha_1, \alpha_2) &= \min_{\alpha_2 \in \Delta A_2} U_1(\alpha_1^{\max}, \alpha_2) \\ &\leq U_1(\alpha_1^{\max}, \alpha_2) \quad \forall \alpha_2 \in \Delta A_2 \end{aligned}$$

and

$$\begin{aligned} \min_{\alpha_2 \in \Delta A_2} \max_{\alpha_1 \in \Delta A_1} U_1(\alpha_1, \alpha_2) &= \max_{\alpha_1 \in \Delta A_1} U_1(\alpha_1, \alpha_2^{\min}) \\ &\geq U_1(\alpha_1, \alpha_2^{\min}) \quad \forall \alpha_1 \in \Delta A_1 \end{aligned}$$

But

$$\begin{aligned}\max_{\alpha_1 \in \Delta A_1} \min_{\alpha_2 \in \Delta A_2} U_1(\alpha_1, \alpha_2) &= \min_{\alpha_2 \in \Delta A_2} \max_{\alpha_1 \in \Delta A_1} U_1(\alpha_1, \alpha_2) \\ &= U_1(\alpha_1^{\max}, \alpha_2^{\min})\end{aligned}$$

implies that

$$U_1(\alpha_1, \alpha_2^{\min}) \leq U_1(\alpha_1^{\max}, \alpha_2^{\min}) \leq U_1(\alpha_1^{\max}, \alpha_2)$$

$\forall \alpha_2 \in \Delta A_2$  and  $\forall \alpha_1 \in \Delta A_1$ .

Hence,  $(\alpha_1^{\max}, \alpha_2^{\min})$  is an equilibrium.

## Interchangeability

If  $\alpha$  and  $\alpha'$  are *NE* in a zero-sum game, then so are  $(\alpha_1, \alpha'_2)$  and  $(\alpha'_1, \alpha_2)$ .

– Since  $\alpha$  and  $\alpha'$  are equilibria

$$U_1(\alpha_1, \alpha_2) \geq U_1(\alpha'_1, \alpha_2) \text{ and } U_2(\alpha'_1, \alpha'_2) \geq U_2(\alpha'_1, \alpha_2),$$

and because  $U_1 = -U_2$

$$U_1(\alpha'_1, \alpha'_2) \leq U_1(\alpha'_1, \alpha_2).$$

Therefore,

$$U_1(\alpha_1, \alpha_2) \geq U_1(\alpha'_1, \alpha_2) \geq U_1(\alpha'_1, \alpha'_2). \quad (1)$$

and similar analysis gives that

$$U_1(\alpha_1, \alpha_2) \leq U_1(\alpha_1, \alpha'_2) \leq U_1(\alpha'_1, \alpha'_2). \quad (2)$$



– (1) and (2) yield

$$U_1(\alpha_1, \alpha_2) = U_1(\alpha'_1, \alpha_2) = U_1(\alpha_1, \alpha'_2) = U_1(\alpha'_1, \alpha'_2)$$

– Since  $\alpha$  is an equilibrium

$$U_2(\alpha_1, \alpha''_2) \leq U_2(\alpha_1, \alpha_2) = U_2(\alpha_1, \alpha'_2)$$

for any  $\alpha''_2 \in \Delta A_2$ , and since  $\alpha'$  is an equilibrium

$$U_1(\alpha''_1, \alpha'_2) \leq U_1(\alpha'_1, \alpha'_2) = U_1(\alpha_1, \alpha'_2)$$

for any  $\alpha''_1 \in \Delta A_1$ . Therefore,  $(\alpha_1, \alpha'_2)$  is an equilibrium and similarly also  $(\alpha_1, \alpha'_2)$ .